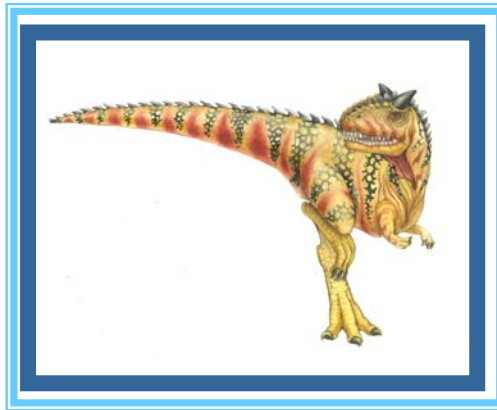


# Chapter 14: Protection

---





# Chapter 14: Protection

---

- Goals of Protection
- Principles of Protection
- Domain of Protection
- Access Matrix
- Implementation of Access Matrix
- Access Control
- Revocation of Access Rights
- Capability-Based Systems
- Language-Based Protection





# Objectives

---

- Discuss the goals and principles of protection in a modern computer system
- Explain how protection domains combined with an access matrix are used to specify the resources a process may access
- Examine capability and language-based protection systems





# Goals of Protection

---

- Operating system consists of a collection of objects, hardware or software
- Each object has a unique name and can be accessed through a well-defined set of operations
- Protection problem - ensure that each object is accessed correctly and only by those processes that are allowed to do so





# Principles of Protection

---

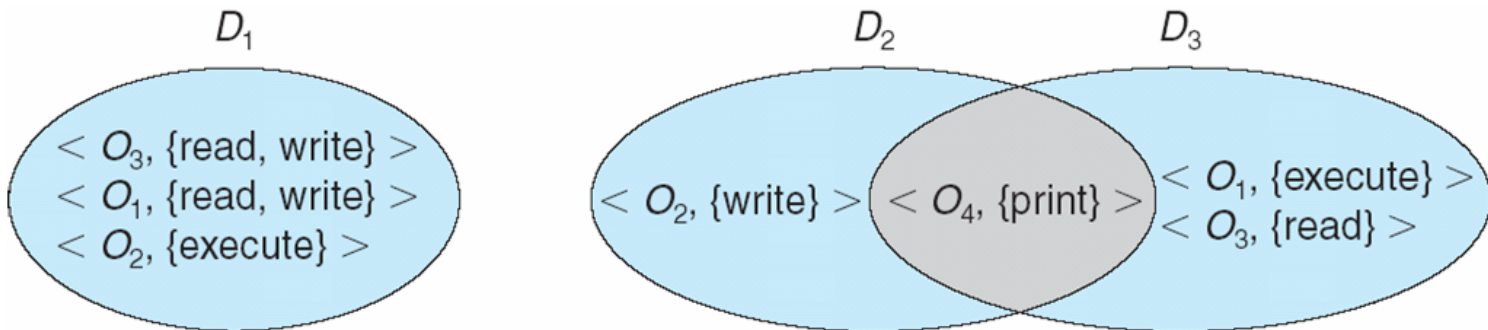
- Guiding principle – principle of least privilege
  - Programs, users and systems should be given just enough privileges to perform their tasks





# Domain Structure

- Access-right =  $\langle \text{object-name}, \text{rights-set} \rangle$   
where *rights-set* is a subset of all valid operations that can be performed on the object.
- Domain = set of access-rights





# Domain Implementation (UNIX)

---

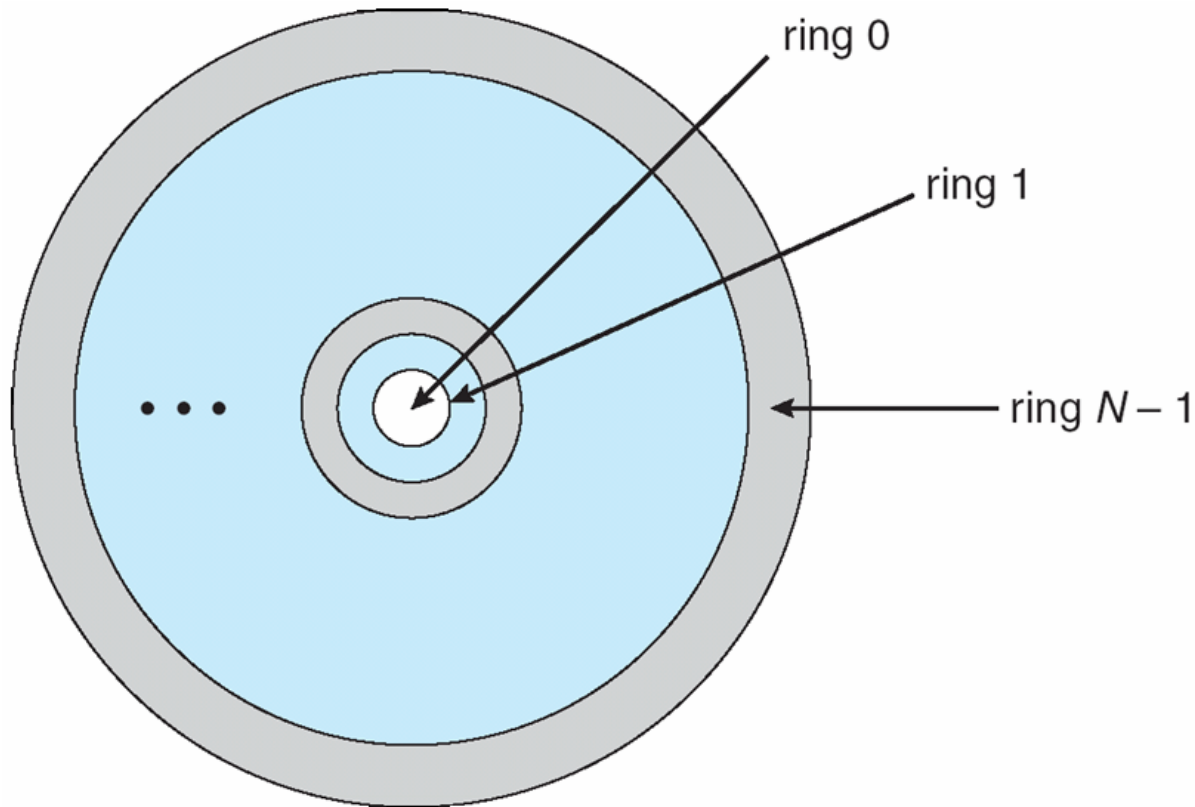
- System consists of 2 domains:
  - User
  - Supervisor
  
- UNIX
  - Domain = user-id
  - Domain switch accomplished via file system
    - ▶ Each file has associated with it a domain bit (setuid bit)
    - ▶ When file is executed and setuid = on, then user-id is set to owner of the file being executed. When execution completes user-id is reset





# Domain Implementation (MULTICS)

- Let  $D_i$  and  $D_j$  be any two domain rings
- If  $j < i \Rightarrow D_i \subseteq D_j$







# Access Matrix

---

- View protection as a matrix (*access matrix*)
- Rows represent domains
- Columns represent objects
- $\text{Access}(i, j)$  is the set of operations that a process executing in Domain <sub>$i$</sub>  can invoke on Object <sub>$j$</sub>





# Access Matrix

object domain	$F_1$	$F_2$	$F_3$	printer
$D_1$	read		read	
$D_2$				print
$D_3$		read	execute	
$D_4$	read write		read write	





# Use of Access Matrix

- If a process in Domain  $D_i$  tries to do “op” on object  $O_j$ , then “op” must be in the access matrix
- Can be expanded to dynamic protection
  - Operations to add, delete access rights
  - Special access rights:
    - ▶ *owner of  $O_i$*
    - ▶ *copy op from  $O_i$  to  $O_j$*
    - ▶ *control –  $D_i$  can modify  $D_j$  access rights*
    - ▶ *transfer – switch from domain  $D_i$  to  $D_j$*





# Use of Access Matrix (Cont)

---

- Access matrix design separates mechanism from policy
  - Mechanism
    - ▶ Operating system provides access-matrix + rules
    - ▶ If ensures that the matrix is only manipulated by authorized agents and that rules are strictly enforced
  - Policy
    - ▶ User dictates policy
    - ▶ Who can access what object and in what mode





# Implementation of Access Matrix

---

- Each column = Access-control list for one object  
Defines who can perform what operation.

Domain 1 = Read, Write  
Domain 2 = Read  
Domain 3 = Read

- Each Row = Capability List (like a key)  
Fore each domain, what operations allowed on what objects.
  - Object 1 – Read
  - Object 4 – Read, Write, Execute
  - Object 5 – Read, Write, Delete, Copy





# Access Matrix of Figure A With Domains as Objects

domain \ object	$F_1$	$F_2$	$F_3$	laser printer	$D_1$	$D_2$	$D_3$	$D_4$
$D_1$	read		read			switch		
$D_2$				print			switch	switch
$D_3$		read	execute					
$D_4$	read write		read write		switch			

**Figure B**





# Access Matrix with Copy Rights

domain \ object	$F_1$	$F_2$	$F_3$
$D_1$	execute		write*
$D_2$	execute	read*	execute
$D_3$	execute		

(a)

domain \ object	$F_1$	$F_2$	$F_3$
$D_1$	execute		write*
$D_2$	execute	read*	execute
$D_3$	execute	read	

(b)





# Access Matrix With *Owner* Rights

<div>object domain</div>	$F_1$	$F_2$	$F_3$
$D_1$	owner execute		write
$D_2$		read* owner	read* owner write
$D_3$	execute		

(a)

<div>object domain</div>	$F_1$	$F_2$	$F_3$
$D_1$	owner execute		write
$D_2$		owner read* write*	read* owner write
$D_3$		write	write

(b)







# Modified Access Matrix of Figure B

object domain	$F_1$	$F_2$	$F_3$	laser printer	$D_1$	$D_2$	$D_3$	$D_4$
$D_1$	read		read			switch		
$D_2$				print			switch	switch control
$D_3$		read	execute					
$D_4$	write		write		switch			





# Access Control

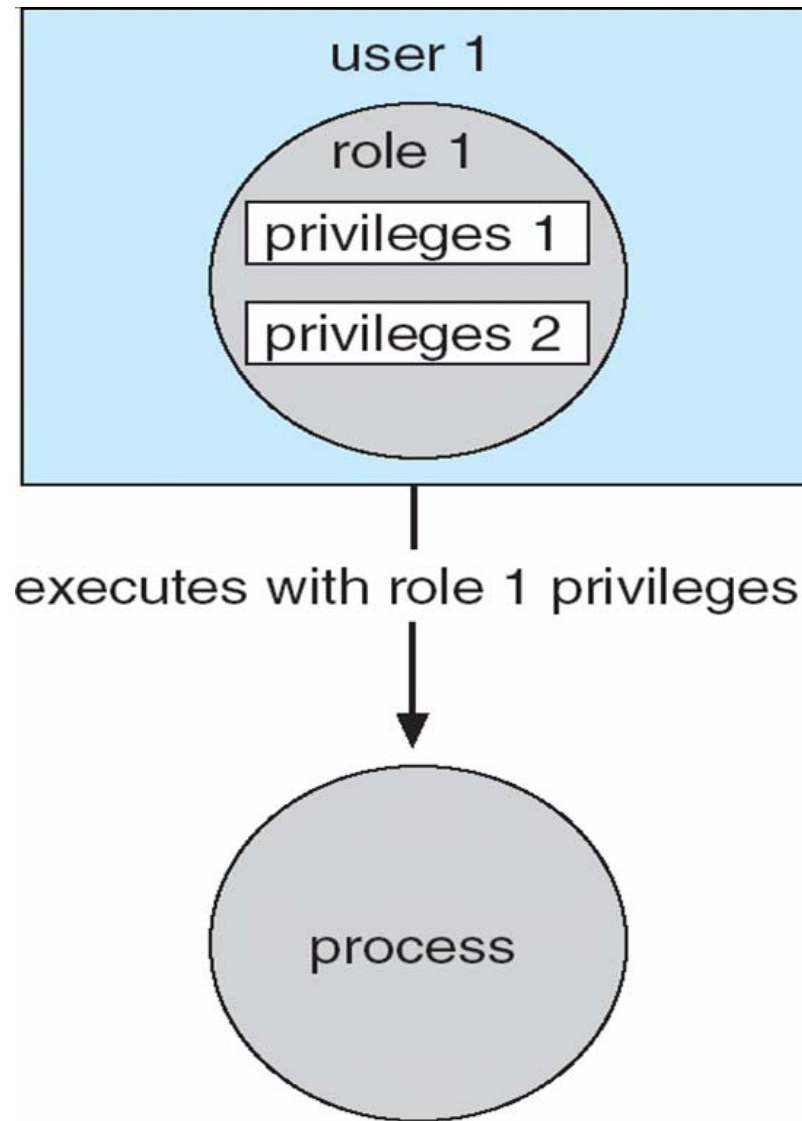
---

- Protection can be applied to non-file resources
- Solaris 10 provides **role-based access control (RBAC)** to implement least privilege
  - Privilege is right to execute system call or use an option within a system call
  - Can be assigned to processes
  - Users assigned roles granting access to privileges and programs





# Role-based Access Control in Solaris 10





# Revocation of Access Rights

---

- **Access List** – Delete access rights from access list
  - Simple
  - Immediate
  
- **Capability List** – Scheme required to locate capability in the system before capability can be revoked
  - Reacquisition
  - Back-pointers
  - Indirection
  - Keys





# Capability-Based Systems

---

- Hydra
  - Fixed set of access rights known to and interpreted by the system
  - Interpretation of user-defined rights performed solely by user's program; system provides access protection for use of these rights
- Cambridge CAP System
  - Data capability - provides standard read, write, execute of individual storage segments associated with object
  - Software capability -interpretation left to the subsystem, through its protected procedures





# Language-Based Protection

---

- Specification of protection in a programming language allows the high-level description of policies for the allocation and use of resources
- Language implementation can provide software for protection enforcement when automatic hardware-supported checking is unavailable
- Interpret protection specifications to generate calls on whatever protection system is provided by the hardware and the operating system





# Protection in Java 2

---

- Protection is handled by the Java Virtual Machine (JVM)
- A class is assigned a protection domain when it is loaded by the JVM
- The protection domain indicates what operations the class can (and cannot) perform
- If a library method is invoked that performs a privileged operation, the stack is inspected to ensure the operation can be performed by the library





# Stack Inspection

protection domain:	untrusted applet	URL loader	networking
socket permission:	none	*.lucent.com:80, connect	any
class:	gui: ... get(url); open(addr); ...	get(URL u): ... doPrivileged { open('proxy.lucent.com:80'); } <request u from proxy> ...	open(Addr a): ... checkPermission (a, connect); connect (a); ...





# End of Chapter 14

---

